A decision tool based on bilevel optimization for the allocation of water resources in a hierarchical system

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Abstract

This paper addresses the optimal allocation of water among competing stakeholders during a finite planning horizon. We focus on those water systems where there are two levels of decision making organized according to a hierarchical framework. At the upper level, a central authority allocates water to demand points having regard to environmental and sustainability issues as well as balancing water users' supply/demand. At the lower level of the hierarchy, demand point managers allocate water to users prioritizing economic strategies. On the other hand, when it comes to allocating limited resources that affect public welfare, the authority in charge can also use different political instruments such as fees to influence the decisions made at those levels of decision making that are not directly within its competence. We propose a multiobjective multifollower bilevel optimization problem that aims to fulfill the central authority goals while including the reaction of the demand point managers in terms of optimization problems as constraints. Using the well-known Karush–Kuhn–Tucker approach, we transform the bilevel model into an equivalent multiobjective mixed-integer single-level model for which we provide tight big- M values. For the purpose of showing the versatility of the model, extensive computational experiments on a set of instances have been carried out. The results show that the optimization problem can be solved to optimality in small computing times using off-the-shelf mixed-integer solvers even for complex water systems and long planning periods. In addition, they illustrate the effect of imposing fees on the achievement of the central authority's objectives.

Keywords: water system management; bilevel optimization; multiobjective; multifollower; exact solution; decision tool

1. Introduction

Distributing water has become a very important issue due to its increasing uses and needs. In cases where the demand for water is small compared to its availability, conflicts do not arise. However, in

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general, water resources are too limited to meet all the needs of water users, and the use of water by one user affects the availability for other users. Moreover, water users usually pursue different and even conflicting objectives. Therefore, water system managers face conflicts when users compete for water and making decisions become increasingly complex. On the other hand, there is an increasing interest in improving water resources management by taking into account sustainability and environmental protection of aquatic ecosystems as well as giving users the opportunity of participating in the decision-making process (Hassing et al., 2009; Ruiz-Villaverde and García-Rubio, 2017). Optimization techniques have provided support to deal with water system management. In the literature, the long-established modeling approach assumes centralized planning, that is, the water distribution decision process is controlled by a single decision maker. In some cases, multiple decision makers are allowed, who agree to collaborate to achieve common goals in a decentralized coordinated planning process. Therefore, optimization models with a single decision level and one or several criteria have been proposed in the literature; see Brown et al. (2015), Calvete and Mateo (1995), Llopis-Albert et al. (2018), Mala-Jetmarova et al. (2017), Roozbahani et al. (2013), Tayfur (2017), Udías et al. (2012), and references therein. Most of these papers, when explicitly taking into account environmental issues besides economic issues, propose a multiobjective model.

However, there are many practical problems in which the agents involved in water allocation are not necessarily willing to cooperate. For some of them, environmental issues must be considered a priority. For others, obtaining the best economic performance is the only thing that matters. In order to deal with this conflict, and bearing in mind that usually a hierarchy can be recognized in the decision-making process with a central authority having a prevailing role, a bilevel optimization approach for the optimal allocation of water is more appropriate, such as that proposed in this paper. This approach allows the central authority the achievement of environmental goals and the use of fees to encourage a balanced allocation of water, while the demand points decide on this allocation. The idea of a central authority influencing decision makers has already been applied in different fields. Among others, Önal et al. (1995) model the distribution of agricultural credits between farm groups aiming to improve the agriculture sector's performance. Amouzegar and Moshirvaziri (1999) deal with managing hazardous waste aiming to maximize social welfare via taxation. Zhao et al. (2013) study how to control water pollution to attain the desired water quality at the lowest environmental cost. Bostian et al. (2015) aim to control nitrogen loading in the watershed to achieve a trade-off between agricultural production and water quality. The review by Sinha et al. (2018) includes several references that deal with the application of bilevel optimization in environmental economics, where an authority or regulator at the upper level of decision making tax those entities at the lower level that are polluting the environment as a result of their operations. In a hierarchized sustainable supply chain network design, Chalmardi and Camacho-Vallejo (2019) propose a bilevel model in which the government acts as the leader to incentivize the use of cleaner technologies by offering financial incentives to the supply chain's managers. For eco-industrial parks, Aviso et al. (2010) propose a bilevel fuzzy optimization model to explore the effect of charging fees for the purchase of freshwater and for the treatment of wastewater in optimizing the water exchange network of plants and Bi et al. (2019) model the distribution of water to minimize overall water consumption.

Bilevel optimization models those problems having a hierarchical framework in which there are two decision levels whose decision makers, besides having different goals, only have control over some of the decisions to be made. The decision maker at the upper level of the hierarchy aims to

optimize his/her objective function under a set of constraints that take into account the reaction of the decision maker at the lower level of the hierarchy to the plan of the leader. Therefore, bilevel programs are formulated as optimization problems, which involve another optimization problem in the constraint set. Regarding the application of bilevel optimization in water-allocation problems, Guo et al. (2012) propose a bilevel model for a multireservoir operation model in an inter-basin water transfer supply project. The upper level controls the distribution of water resources among water exporting and importing regions using a set of water-transfer rules. The individual reservoir manager, at the lower level of the hierarchy, controls the water-supply process by hedging rules. The authors develop an improved particle swarm optimization that is used to solve both upper and lower level models transferring the value of the corresponding variables from one model to another, and apply it to a three-reservoir system in a province of Northeast China. Zhu et al. (2017) consider the problem of transferring and supplying water and present a bilevel model, which is solved using an adaptive genetic algorithm. The upper level's goal is to minimize the actual annual average transferred water while the lower level aims to minimize water shortage for all users. A fuzzy approach is often taken to solve the bilevel problem when the upper and lower decision makers are able to cooperate with each other. Under this hypothesis, the bilevel problem can usually be solved sequentially and thus it is simpler to handle. This approach is adopted by Lv et al. (2010) and Xu et al. (2012) for planning water resources management systems. The aim is to balance the degree of satisfaction between the upper and lower levels. Chen et al. (2017) propose a bilevel model in which the upper level decision maker aims to minimize the discharge of pollutants when deciding the water-allocation strategy, whereas the lower level representing the regional authority maximizes the economic benefits. To solve the problem, an interactive algorithm is developed based on the concepts of satisfaction and tolerance membership functions of the fuzzy theory.

The contribution of this paper is to propose a model for water resource allocation in a general complex noncooperative hierarchical water system. The model is flexible enough to allow consideration of very different water system configurations, the evolution over time of the system, and environmental and economic issues. In addition, the model enables the use of fees that can be charged by the decision maker at the highest level of the hierarchy to influence the decisions made at those levels of decision making that are not directly within its competence. The result is a decision tool that can be used to evaluate different strategies in the optimal allocation of the available water when the reaction of the users is included in the decision process.

The study considers the management during a finite planning horizon, divided into time periods, of a water resources system consisting of rivers, reservoirs, distribution channels, and water demand points, each having several water users. Due to water shortages, conflicts arise among the demand points. The aim is to determine how to allocate the available water taking into account the hierarchical structure of the decision process. On the one hand, a central authority representing the government or water system authorities decides on the global amount of water allocated to each demand point. It aims to distribute water in accordance with environmental aspects and the overall satisfaction of user demand, having the possibility of charging fees on water allocated to users. On the other hand, managers at the demand points (which represent water users' communities) decide on the different uses of the allocated water, that is, on the distribution to their water users based on maximizing the net economic return. Both decision levels act under a hierarchical structure. This means that the central authority at the upper level of the decision process has control of the water

distribution to the demand points, but the reaction of the managers who decide on water uses is taken into account in the model as a constraint.

This issue is modeled as a multiobjective lexicographic linear bilevel optimization problem with a leader and several followers, one for each demand point. First, taking into account that each follower problem only involves its own variables and the variables controlled by the upper level, this model is transformed into a multiobjective lexicographic linear bilevel problem with a single follower. Then, to solve the problem, we propose a procedure that takes advantage of the fact that, after the above-mentioned transformation, the lower level problem is a linear one. Hence, Karush–Kuhn–Tucker (KKT) conditions are necessary and sufficient for optimality, and the multiobjective lexicographic bilevel problem can be transformed into an equivalent multiobjective lexicographic mixed-integer linear single-level problem by introducing additional binary variables. This linearization process involves the use of big- M constants for which we provide tight valid values. The resulting optimization problem can be solved to optimality using off-the-shelf mixed-integer solvers. It is worth pointing out that the model can efficiently handle complex water systems with long planning periods. To illustrate the performance of this approach, we have generated a set of instances that handle a variety of water system configurations under different water availability scenarios. Extensive computational experiments on this set show the versatility of the model as well as the short computing times required to solve it. As a result, this work provides a mathematical model that can be used as a decision tool by the central authority to assess the degree of fulfillment of its objectives and to quickly evaluate the consequences of different simulated scenarios when deciding water allocation.

The rest of the paper is structured as follows. Following the Introduction section, Section 2 provides a brief summary of bilevel optimization. Section 3 describes in detail the formulation of the bilevel model and the procedure used to solve it. Section 4 presents the results of the computational experiments carried out. Finally, Section 5 concludes the paper.

2. Background on bilevel optimization problems

As mentioned above, bilevel optimization has been proposed to deal with hierarchical decision processes with two levels of decision. When there is a leader and a single follower, the bilevel optimization model can be formulated as

$$\begin{aligned}
 & \min_x && F(x, y) \\
 & \text{subject to} && \\
 & && G_j(x, y) \leq 0, \quad j = 1, \dots, q, \\
 & \text{where, for every } x \text{ fixed, } y \text{ solves} && \\
 & \min_y && f(x, y) \\
 & \text{subject to} && \\
 & && g_h(x, y) \leq 0, \quad h = 1, \dots, p,
 \end{aligned} \tag{1}$$

where $x \in \mathbb{R}^n$ are the upper level variables controlled by the leader, and $y \in \mathbb{R}^m$ are the lower level variables controlled by the follower; $F, f : \mathbb{R}^{n+m} \rightarrow \mathbb{R}$ are the upper and lower level objective

functions, respectively; and $G_j(x, y) \leq 0$, $g_h(x, y) \leq 0$ refer, respectively, to constraints of the upper and lower levels.

The bilevel problem is the leader's problem, who must anticipate the follower's reaction when looking for his/her best decision. The follower is free to optimize his/her objective function once the leader sets the value of the upper level variables. Therefore, for a given x , the follower solves the lower level problem:

$$\begin{aligned} \min_y \quad & f(x, y) \\ \text{subject to} \quad & g_h(x, y) \leq 0, \quad h = 1, \dots, p. \end{aligned} \quad (2)$$

Let $V = \{(x, y) : G_j(x, y) \leq 0, j = 1, \dots, q, g_h(x, y) \leq 0, h = 1, \dots, p\}$. Let $M(x)$ be the optimal solution set of problem (2). Let $T = \{(x, y) : (x, y) \in V, y \in M(x)\}$. A point $(x, y) \in T$ is a so-called bilevel feasible solution, that is, a feasible solution of problem (1). A point $x \in \mathbb{R}^n$ is called permissible if a point $y \in \mathbb{R}^m$ exists so that $(x, y) \in T$. Hence, the bilevel problem (1) can be equivalently formulated as

$$\begin{aligned} \min_x \quad & F(x, y) \\ \text{s.t.} \quad & (x, y) \in T. \end{aligned}$$

The bilevel problem is nonconvex and difficult to deal with and solve. In fact, complications arise when there are multiple optima in the lower level problem, that is to say, $M(x)$ is not a singleton for some permissible x . If the upper level objective function is sensitive to different values of $y \in M(x)$, it is necessary to give a rule to select $y^* \in M(x)$ in order to evaluate F . Several assumptions have been proposed in the literature, the most common being the optimistic approach, which assumes that the upper level decision maker has the right to influence the lower level decision maker so that the latter selects y^* to provide the best value of F . In this case, the upper level objective function is minimized over x and y . This is the approach taken in this paper. It is also worth noting that even when all the functions involved are linear, the bilevel problem is strongly NP-hard. Calvete et al. (2012) provide an overall view of the main difficulties that arise when dealing with bilevel problems and distinguish them from classical single-level optimization problems. Bard (1998) and Dempe (2002) are textbooks that deal with most of the theoretical issues on bilevel optimization. Colson et al. (2007), Dempe (2018), Sinha et al. (2018), and Dempe and Zemkoho (2020) provide recent surveys that cover applications as well as major theoretical developments.

3. A multiobjective multifollower bilevel optimization model for water system management

In this section, we mathematically describe the problem addressed by the central authority when allocating water and formulate the bilevel problem that models it. Water enters the system at various locations and is distributed to demand points or stored in reservoirs; otherwise, it leaves the system. From the demand points, the water is allocated to water users aiming to satisfy their demands. To take into account that demand and inflows change over time, we consider the evolution of the water

system during a finite planning horizon divided into periods. For the water system management, it is necessary to know not only if there is enough water overall to satisfy demand but if water is available for allocation during the required period. We assume that in each period of the planning horizon, the inflows to the system and every water user demand are known or can be properly estimated. Moreover, minimum downstream requirements and minimum storage in reservoirs are also known and must be prioritized when determining the optimal water distribution to meet environmental and sustainability issues. The goal is to determine the best way of allocating available water. The central authority, playing the role of leader, controls the water that is allocated to the demand points, as well as minimum needs and the flows leaving the system. Each demand point, acting as a follower, controls the allocation of water to its water users.

3.1. Assumptions, notations, and variables

Let T denote the set of periods of the planning horizon. Let R be the set of reservoirs, W be the set of demand points, and K be the set of water users. Each reservoir $r \in R$ is connected to other reservoirs and/or to demand points. Let $R_r \subseteq R$ and $W_r \subseteq W$ be, respectively, the sets of reservoirs and demand points to which r is connected. This means that water can flow from the reservoir r to any node in $R_r \cup W_r$. Each demand point $w \in W$ is only connected to the set $K_w \subseteq K$ of water users to which w can allocate water. We assume that the sets of water users associated with different demand points are disjoint, that is, $K_w \cap K_{\tilde{w}} = \emptyset$ if $w \neq \tilde{w}$. Hence, $\{K_w\}_{w \in W}$ constitute a partition of K .

Let I_r^t be the inflow in reservoir $r \in R$ during period $t \in T$. For each water user $k \in K_w$, $w \in W$, let D_{wk}^t denote the total demand during period $t \in T$ and b_{wk}^t be the economic return obtained by the demand point w per unit of water allocated to the user k during period t . Let P_{wk}^t , $w \in W$, $k \in K_w$, $t \in T$, be an upper bound on the price, which the central authority can charge the demand point w for the allocation of water to the water user k during period t . We assume that $P_{wk}^t \leq b_{wk}^t$, $w \in W$, $k \in K_w$, $t \in T$, that is, the prices are never greater than the economic return of the corresponding water user.

Each reservoir $r \in R$ has a capacity and a minimum level of storage to be guaranteed at each period as required by the central authority. We denote x_{rr}^0 the amount of water available at the reservoir $r \in R$ at the beginning of the planning period. Moreover, the links between reservoirs and demand points represent either a river or a channel. Hence, in general, a link representing a river has a lower bound meaning a minimum downstream requirement and does not have an upper bound. This also applies to the water exiting the system. Conversely, a channel has an upper bound referring to its capacity and a lower bound equal to zero. We denote the lower bounds by L and the upper bounds by U , with the proper indices in each case. For example, L_{rr} and U_{rr} denote, respectively, the minimum level of storage and the capacity of reservoir $r \in R$, while L_{rw} and U_{rw} are, respectively, zero and the capacity of the channel if there is a channel connecting the reservoir $r \in R$ and the demand point $w \in W$.

It is worth noting that, for a particular planning period, it cannot be guaranteed *a priori* that enough water will be available to satisfy all the needs. In fact, even if overall enough water was available during the planning horizon, depending on the characteristics of the water system, it is possible that not enough water would be available to ensure even the minimum level of storage and

the minimum downstream requirements in some periods of the planning horizon. Also, water users can request more water than available in some periods. Therefore, to enforce minimum requirements and to maintain potential demands as mandatory demands that must be supplied would lead to an infeasible problem. However, in these cases where there is not enough water, it is even more important for the central authority to know how to distribute water according to its goals. Hence, the possibility of deficits is considered.

In order to formulate the multiobjective lexicographic multifollower bilevel optimization problem (MLMF-BP), we define the following upper level decision variables, which are controlled by the central authority:

$x_{rr}^t, r \in R, t \in T:$	storage at reservoir r at the end of period t
$x_{r\tilde{r}}^t, r \in R, \tilde{r} \in R_r, t \in T:$	flow from reservoir r to reservoir $\tilde{r}$ during period t
$x_{rw}^t, r \in R, w \in W_r, t \in T:$	water sent from reservoir r to demand point w during period t
$x_r^t, r \in R, t \in T:$	water leaving the system from reservoir r during period t
$s_{rr}^t, r \in R, t \in T:$	storage deficit at reservoir r at the end of period t
$s_{r\tilde{r}}^t, r \in R, \tilde{r} \in R_r, t \in T:$	minimum downstream requirement deficit from reservoir r to reservoir $\tilde{r}$ during period t
$s_r^t, r \in R, t \in T:$	minimum downstream requirement deficit for the water exiting the system from reservoir r during period t
$p_{wk}^t, w \in W, k \in K_w, t \in T:$	price charged to the demand point w per unit of water allocated to the water user k during period t

Note that the variables x_r^t and s_r^t are only defined for those reservoirs r from which the water can leave the system. Also, it must be noted that the aim of the central authority when using the variables p_{wk}^t is not to maximize the return achieved from these prices but to influence the water allocation of demand points.

The lower level decision variables controlled by the demand point $w \in W$ are defined as follows:

$$y_{wk}^t, k \in K_w, t \in T: \text{ water allocated to the water user } k \text{ during period } t.$$

As an illustration, a simple water system is shown at the top of Fig. 1. This consists of a sequential water system with two reservoirs, three demand points, and seven users. To help understand the evolution of the system over time and the variables involved, at the bottom of the same figure we have represented a diagram of the system over two periods. The variables associated to each arc are shown next to the arc.

3.2. The multiobjective lexicographic multifollower bilevel optimization model

The MLMF-BP can be formulated as follows:

$$\min_{x,s,p,y} f_1^1 = \sum_{t \in T} \sum_{r \in R} \left(s_{rr}^t + s_r^t + \sum_{\tilde{r} \in R_r} s_{r\tilde{r}}^t \right) \quad (3a)$$

Fig. 1. A simple water system on top together with a representation of the flows and the variables involved when there are two periods of time, $T = \{1, 2\}$. In this case, $R = \{r_1, r_2\}$, $W = \{w_1, w_2, w_3\}$, $K = \{k_1, \dots, k_7\}$, $R_{r_1} = \{r_2\}$, $W_{r_1} = \{w_1, w_2\}$, and $K_{w_1} = \{k_1, k_2\}$.

$$\max_{x,s,p,y} f_1^2 = \sum_{t \in T} \sum_{w \in W} \sum_{k \in K_w} \frac{y_{wk}^t}{D_{wk}^t} \quad (3b)$$

$$\min_{x,s,p,y} f_1^3 = \sum_{t \in T} \sum_{w \in W} \sum_{k \in K_w} p_{wk}^t \quad (3c)$$

s.t.

$$x_{rr}^{t-1} + I_r^t + \sum_{\tilde{r}: r \in R_{\tilde{r}}} x_{\tilde{r}r}^t = x_{rr}^t + x_r^t + \sum_{\tilde{r} \in R_r} x_{r\tilde{r}}^t + \sum_{w \in W_r} x_{rw}^t, \quad r \in R, t \in T \quad (3d)$$

$$x_{rr}^t + s_{rr}^t \geq L_{rr} \quad r \in R, t \in T \quad (3e)$$

$$x_r^t + s_r^t \geq L_r \quad r \in R, t \in T \quad (3f)$$

$$x_{r\tilde{r}}^t + s_{r\tilde{r}}^t \geq L_{r\tilde{r}} \quad r \in R, \tilde{r} \in R_r, t \in T \quad (3g)$$

$$x_{rr}^t \leq U_{rr} \quad r \in R, t \in T \quad (3h)$$

$$x_{r\tilde{r}}^t \leq U_{r\tilde{r}} \quad r \in R, \tilde{r} \in R_r, t \in T \quad (3i)$$

$$x_{rw}^t \leq U_{rw} \quad r \in R, w \in W_r, t \in T \quad (3j)$$

$$p_{wk}^t \leq P_{wk}^t \quad w \in W, k \in K_w, t \in T \quad (3k)$$

$$x_r^t, x_{rr}^t, x_{r\tilde{r}}^t, x_{rw}^t, s_r^t, s_{rr}^t, s_{r\tilde{r}}^t \geq 0 \quad r \in R, \tilde{r} \in R_r, w \in W_r, t \in T \quad (3l)$$

$$p_{wk}^t \geq 0 \quad w \in W, k \in K_w, t \in T, \quad (3m)$$

where, for each demand point $w \in W$, the variables $y_{wk}^t, k \in K_w, t \in T$, solve:

$$\max_{y_w} f_{2w} = \sum_{t \in T} \sum_{k \in K_w} (b_{wk}^t - p_{wk}^t) y_{wk}^t \quad (3n)$$

s.t.

$$\sum_{k \in K_w} y_{wk}^t \leq \sum_{r \in R: w \in W_r} x_{rw}^t, \quad t \in T \quad (3o)$$

$$y_{wk}^t \leq D_{wk}^t, \quad k \in K_w, t \in T \quad (3p)$$

$$y_{wk}^t \geq 0, \quad k \in K_w, t \in T. \quad (3q)$$

The objective functions (3a)–(3c) lexicographically optimize the three ranked objectives of the central authority. Lexicographic optimality (Ehrgott, 2005) implies a ranking of the objectives. That is to say, it is assumed that the first objective is more important by far than the second objective, which in turn is more important by far than the third objective, and so on. Thus, the optimization

of an objective is not considered until the optimality of the objectives, which precede it in importance, has been established. Concerning the MLMF-BP, the goal with the highest priority (3a) is to provide water to satisfy environmental and sustainability issues, that is, to meet the minimum requirements (minimum storage in reservoirs and minimum downstream requirements). Hence, the total deficit is minimized. If the central authority has to differentiate among those needs in case of not having enough water available to satisfy all of them, an additional weight coefficient can be associated with the deficit variables. This coefficient would allow us to emphasize the relevance of each need.

The second prioritized goal (3b) refers to the satisfaction of water users. In this paper, the concept of satisfaction proposed by Babel et al. (2005) as the ratio of water allocated over the demand is used. In addition, we propose to consider the utilitarian approach in which the goal is to maximize the total satisfaction of demand users measured by the sum of the individual satisfaction of each of them. As mentioned above, an additional weight coefficient can be associated with the individual satisfaction to emphasize its relevance. Other approaches that propose different strategies to deal with the allocation of resources across multiple users can be seen in Karsu and Erkan (2020) and references therein.

As mentioned above, the aim of the central authority in setting the prices is not to collect as much as possible but to be able to influence the allocation of water to the water users. Hence, from the, generally large, set of prices that provide the outcome sought by the central authority, we propose to select a set of low prices, and so a third prioritized goal (3c) is introduced which minimizes the sum of the prices. This will allow the demand points to be charged prices that together add up to as little as possible, always ensuring that the overall satisfaction represented by the second objective is maximized. In order to keep the model as simple as possible, within the complexity that is inherent in bilevel models, we have assumed that all the objective functions should be linear. Thus, the third objective function minimizes the total sum of prices instead of the total paid by users. The function $\sum_{t \in T} \sum_{w \in W} \sum_{k \in K_w} p_{wk}^t y_{wk}^t$ is not linear and we think that f_1^3 defined in (3c) fairly well captures the aim of the central authority. Nevertheless, it is important to note that other objective functions can be considered, which gives versatility to the proposed model. For instance, the central authority may want to minimize the highest price.

To continue with the description of the model, constraints (3d) ensure the conservation of flow. Constraints (3e) refer to the minimum level of storage in the reservoirs. Constraints (3f) refer to the minimum downstream requirements of the water leaving the water system. Constraints (3g) refer to the minimum downstream requirements of the link connecting r and $\tilde{r}$. Constraints (3h) ensure that the capacity of reservoirs is not exceeded. Constraints (3i) and (3j) allow us to guarantee that the capacity of each channel is not exceeded if the link corresponds to a channel, otherwise the upper bound is infinite. Constraints (3k) establish upper bounds on the prices p_{wk}^t , which seems appropriate to avoid unreasonable values for the prices. These upper bounds could be expressed, for instance, as a percentage of the economic return. Other price-related constraints could also be included in the model such as equal prices for certain water users even if they depend on different demand points. Constraints (3l) and (3m) guarantee that all variables controlled by the central authority are nonnegative.

The lower level problem associated with the demand point $w \in W$ is defined by (3n)–(3q). After knowing the amount of water available, the manager of each demand point distributes the water to maximize the net economic return (revenues less costs) from water allocated to water users (3n).

Constraints (3o) ensure that at most the available water at demand point $w \in W$ is distributed among its water users $k \in K_w$. Constraints (3p) guarantee that no water user receives more than he/she demands. Constraints (3q) establish the nonnegativity of the lower level variables.

Problem (3) is an MLMF-BP with $|W|$ followers, where $|W|$ stands for the cardinality of W . Moreover, each follower problem involves only its own variables and the upper level variables. Hence, in accordance with the definition introduced in Calvete and Galé (2007), the followers are independent. By directly extending Theorem 3.1 from that paper to problem (3), this problem is equivalent to a multiobjective lexicographic bilevel optimization problem with one follower. The new lower level problem is obtained by considering the sum of the objective functions f_{2w} of the $|W|$ lower level problems as the objective function, and putting together all the lower level constraints. Consequently, problem (3) can be reformulated as follows:

$$\text{lex } [\min f_1^1, \max f_1^2, \min f_1^3] \quad (4a)$$

s.t.

$$(3d)–(3m), \quad (4b)$$

where the variables y_{wk}^t , $w \in W, k \in K_w, t \in T$, solve

$$\max_y f_2 = \sum_{w \in W} \sum_{t \in T} \sum_{k \in K_w} (b_{wk}^t - p_{wk}^t) y_{wk}^t \quad (4c)$$

s.t.

$$\sum_{k \in K_w} y_{wk}^t \leq \sum_{r \in R: w \in W_r} x_{rw}^t, \quad w \in W, t \in T \quad (4d)$$

$$y_{wk}^t \leq D_{wk}^t, \quad w \in W, k \in K_w, t \in T \quad (4e)$$

$$y_{wk}^t \geq 0, \quad w \in W, k \in K_w, t \in T, \quad (4f)$$

where lex means to lexicographically optimize the three objectives.

3.3. Solving the MLMF-BP

In this section, we propose to reformulate problem (4) as a single-level optimization problem using the KKT conditions of the lower level problem. For linear optimization, KKT conditions are necessary and sufficient for optimality. Hence, they can substitute the lower level problem (4c)–(4f) providing an equivalent problem.

Let us denote by π_w^t , $w \in W, t \in T$, and δ_{wk}^t , $w \in W, k \in K_w, t \in T$, the dual variables associated with the constraints (4d) and (4e), respectively. Then, the dual problem of (4c)–(4f) is

$$\min_{\pi, \delta} \sum_{t \in T} \sum_{w \in W} \left(\sum_{r \in R: w \in W_r} x_{rw}^t \right) \pi_w^t + \sum_{t \in T} \sum_{w \in W} \sum_{k \in K_w} D_{wk}^t \delta_{wk}^t \quad (5a)$$

s.t.

$$\pi_w^t + \delta_{wk}^t \geq b_{wk}^t - p_{wk}^t, \quad w \in W, k \in K_w, t \in T \quad (5b)$$

$$\pi_w^t \geq 0, \quad w \in W, t \in T \quad (5c)$$

$$\delta_{wk}^t \geq 0, \quad w \in W, k \in K_w, t \in T. \quad (5d)$$

Therefore, in addition to the feasibility of primal and dual solutions (constraints (4d)–(4f), (5b)–(5d), the KKT conditions are

$$(\pi_w^t + \delta_{wk}^t - b_{wk}^t + p_{wk}^t)y_{wk}^t = 0, \quad w \in W, k \in K_w, t \in T \quad (6a)$$

$$\left(\sum_{r \in R: w \in W_r} x_{rw}^t - \sum_{k \in K_w} y_{wk}^t \right) \pi_w^t = 0, \quad w \in W, t \in T \quad (6b)$$

$$(D_{wk}^t - y_{wk}^t)\delta_{wk}^t = 0, \quad w \in W, k \in K_w, t \in T. \quad (6c)$$

Nonlinear constraints (6a)–(6c) can be linearized by introducing binary variables v_{wk}^t , v_w^t , and $\tilde{v}_{wk}^t$ as

$$y_{wk}^t \leq M_{wk}^t v_{wk}^t, \quad w \in W, k \in K_w, t \in T \quad (7a)$$

$$\pi_w^t + \delta_{wk}^t - b_{wk}^t + p_{wk}^t \leq N_{wk}^t (1 - v_{wk}^t), \quad w \in W, k \in K_w, t \in T \quad (7b)$$

$$\pi_w^t \leq M_w^t v_w^t, \quad w \in W, t \in T \quad (7c)$$

$$\sum_{r \in R: w \in W_r} x_{rw}^t - \sum_{k \in K_w} y_{wk}^t \leq N_w^t (1 - v_w^t), \quad w \in W, t \in T \quad (7d)$$

$$\delta_{wk}^t \leq \tilde{M}_{wk}^t \tilde{v}_{wk}^t, \quad w \in W, k \in K_w, t \in T \quad (7e)$$

$$D_{wk}^t - y_{wk}^t \leq \tilde{N}_{wk}^t (1 - \tilde{v}_{wk}^t), \quad w \in W, k \in K_w, t \in T \quad (7f)$$

$$v_{wk}^t, \tilde{v}_{wk}^t \in \{0, 1\}, \quad w \in W, k \in K_w, t \in T \quad (7g)$$

$$v_w^t \in \{0, 1\}, \quad w \in W, t \in T, \quad (7h)$$

where M_{wk}^t , N_{wk}^t , M_w^t , N_w^t , $\tilde{M}_{wk}^t$, $\tilde{N}_{wk}^t$ are big enough constants.

Substituting the lower level problem in (4) by its KKT conditions, we obtain the equivalent single-level multiobjective lexicographic mixed-integer problem:

$$\begin{aligned} & \text{lex } [\min f_1^1, \max f_1^2, \min f_1^3] \\ & \text{s.t.} \\ & \quad (3d)–(3m), (4d)–(4f), (5b)–(5d), (7a)–(7h). \end{aligned} \quad (8)$$

Issues related to the choice of appropriate values of these big enough constants have been dealt with in Kleinert et al. (2020) and Pineda and Morales (2019). As cited in Kleinert et al. (2020), our results strongly indicate that problem-specific bounds on the lower level's dual variables need to be investigated if the given bilevel problem is going to be solved using the KKT approach combined with the classical big- M linearization of KKT complementarity conditions. Taking into account the characteristics of the primal and dual problems, the following theorem allows us to derive tight upper bounds.

Theorem 1. For problem (8), valid constants are

$$\begin{aligned} \text{constraints (7a): } & M_{wk}^t = D_{wk}^t \\ \text{constraints (7b): } & N_{wk}^t = b_{wk}^t + \max_{\tilde{k} \in K_w} \{b_{w\tilde{k}}^t\} \\ \text{constraints (7c): } & M_w^t = \max_{\tilde{k} \in K_w} \{b_{w\tilde{k}}^t\} \\ \text{constraints (7d): } & N_w^t = \sum_{k \in K_w} D_{wk}^t \\ \text{constraints (7e): } & \tilde{M}_{wk}^t = b_{wk}^t \\ \text{constraints (7f): } & \tilde{N}_{wk}^t = D_{wk}^t \end{aligned}$$

Proof. According to constraints (4e), $y_{wk}^t \leq D_{wk}^t$. Hence $M_{wk}^t = D_{wk}^t$, $\tilde{N}_w^t = D_{wk}^t$ are valid constants for constraints (7a) and (7f), respectively.

Problem (5) is a minimization problem with nonnegative variables and nonnegative objective function coefficients. Moreover, the right-hand side of each constraint (5b) is at most b_{wk}^t since $p_{wk}^t \geq 0$. Hence, in the optimal solution of problem (5), $\delta_{wk}^t \leq b_{wk}^t$, for all $w \in W$, $k \in K_w$, $t \in T$, whatever the value of the variables π_w^t which guarantee a feasible solution. Note that from each feasible solution having $\delta_{wk}^t > b_{wk}^t$ for some index values, a feasible solution with a lower value of the objective function (5a) can be obtained by making $\delta_{wk}^t = b_{wk}^t$. Hence, $\tilde{M}_{wk}^t = b_{wk}^t$ is a valid constant for constraints (7e).

On the other hand, for each $w \in W$, $t \in T$, the variable $\pi_w^t \leq \max_{\tilde{k} \in K_w} \{b_{w\tilde{k}}^t\}$ in the optimal solution of problem (5). Otherwise, using the same argument as above, $\pi_w^t = \max_{\tilde{k} \in K_w} \{b_{w\tilde{k}}^t\}$ and the corresponding values of δ_{wk}^t would provide a feasible solution with a lower value of the objective function (5a). Hence, $M_w^t = \max_{\tilde{k} \in K_w} \{b_{w\tilde{k}}^t\}$ is a valid constant for constraints (7c).

As a consequence of previous bounds on δ_{wk}^t and π_w^t and taking into account that $p_{wk}^t \leq P_{wk}^t \leq b_{wk}^t$, we can conclude that $N_{wk}^t = b_{wk}^t + \max_{\tilde{k} \in K_w} \{b_{w\tilde{k}}^t\}$ is a valid constant for constraints (7b).

Finally, considering the optimal solution of problem (3), it is always possible to find an optimal solution in which no demand point receives more water than it can allocate to its water users. Thus, for each $w \in W$, $t \in T$, we have $\sum_{r \in R: w \in W_r} x_{rw}^t \leq \sum_{k \in K_w} D_{wk}^t$. Hence, $N_w^t = \sum_{k \in K_w} D_{wk}^t$ is a valid constant for constraints (7d). $\square$

4. Computational experiments

The aim of this section is to present and discuss the computational experiments carried out to illustrate the differences between the optimal water-allocation strategies depending on the water

available and the prices charged, and to show the efficiency of the procedure to solve the MLMF-BP. The numerical experiments have been performed on a PC Intel Core i7-6700 with 3.4 GHz, 32.0 GB of RAM, and Windows 10 64-bit as the operating system. We have solved problem (8) using IBM ILOG CPLEX 12.9.0 with the default settings and the specific tools for multiobjective optimization. The CPLEX stopping criterion has been always set at 7200 seconds.

In addition to solving the MLMF-BP, in these experiments we have also solved what is called the relaxed problem in bilevel optimization. This is a single-level model in which the central authority controls every allocation of water, that is, demand point managers are not taken into account in the decision process. It can be formulated as

$$\begin{aligned} &\text{lex } [\min f_1^1, \max f_1^2] \\ &\text{s.t.} \\ &\quad (3d)–(3j), (3l), (4d)–(4f). \end{aligned} \tag{9}$$

The third objective function f_1^3 and the constraints (3k) and (3m) have been omitted because they force all the prices p_{wk}^t to be zero. Needless to say, when demand points are not accounted for, it does not make sense to charge prices.

On the other hand, the optimal value of f_1^1 in the bilevel problem (3) coincides with the optimal value of the same function in the relaxed problem (9) (both aim to minimize the total deficit, which is controlled by the upper level decision maker, and do not take into account how the demand points are served). Hence, when comparing problems (3) and (9) the main interest lies in the comparison of the upper level objective function f_1^2 . Problem (9) provides a lower bound on f_1^2 of the MLMF-BP, but it does not reflect what actually happens if there exists a hierarchical structure with two levels of decision. Note that once the demand point managers know the value of the upper level variables provided by the optimal solution of the relaxed problem, they assume their role and distribute the allocated water according to their goal. The bilevel feasible solution thus obtained (from now on called the suboptimal solution) is not necessarily optimal for the MLMF-BP. The use of prices can help the central authority to achieve its goals, obtaining objective function values closer to the optimal values of problem (9).

4.1. Set of instances

Most papers in the literature consider cases of study corresponding to small water systems with only one or two reservoirs and they do not provide raw data. When some data are available, their sizes depend largely on the system under consideration. Hence, to have a wide range of scenarios, which allow us to check the versatility of the model as well as to assess the computational time required to solve the model, we have generated six water systems that represent either water systems usually presented in the literature or more complex systems that combine them. They are shown in Fig. 2. Instances WS1 and WS2 consist of a single river, instances WS3 and WS4 consist of a main river with two tributaries, and instances WS5 and WS6 consist of two main rivers connected, respectively, by one and two channels. The number of reservoirs ranges from 4 to 7, the number of demand points ranges from 7 to 12, and the number of water users ranges from 74 to 142. For each instance, we have considered three planning horizons, with 12, 52, and 365 periods.

Fig. 2. Schemes of the water systems of the instances. The numbers in brackets refer to number of water users associated to the corresponding demand point.

Table 1

Randomly generated inflows for $T = 12$

	1	2	3	4	5	6	7	8	9	10	11	12
$(I_r^t)_1$	29,220	25,775	19,998	13,243	7337	3403	3680	6633	12,401	19,018	25,269	29,050
$(I_r^t)_2$	21,841	19,501	15,002	10,015	5312	3069	2577	4801	9259	14,237	18,914	21,947
$(I_r^t)_3$	14,696	12,816	10,083	6669	3666	1911	1856	3228	6167	9730	12,491	14,724

Three yearly inflows (365 periods) have been generated from the following expressions:

$$\begin{aligned}
 (I_r^t)_1 &= 450 \cos \frac{2\pi t}{365} + \text{random}(500, 600), \quad t = 1, \dots, 365 \\
 (I_r^t)_2 &= 0.75 \times (450 \cos \frac{2\pi t}{365} + \text{random}(500, 600)), \quad t = 1, \dots, 365 \\
 (I_r^t)_3 &= 0.5 \times (450 \cos \frac{2\pi t}{365} + \text{random}(500, 600)), \quad t = 1, \dots, 365.
 \end{aligned} \tag{10}$$

Most streams display annual variation due to seasonal changes. In the Mediterranean countries the lowest flows often occur near the end of the spring and during the summer. These sinusoidal functions can represent this behavior.

From the values generated by expression (10), when 12 (respectively, 52) periods are considered, the first 360 (364) values are added up in groups of 30 (7) data each. Table 1 shows the inflow values when 12 periods are considered. Note that less water is available in the central periods. Moreover, distinct scenarios of flow are obtained by multiplying randomly generated inflows by a coefficient β which ranges from 0.1 to 1 in order to simulate variations from a major drought scenario to a high water scenario. For $\beta = 0.8, 0.9$, and 1 the same results are obtained as for $\beta = 0.7$, since there is enough water available to satisfy total demand. Thus, these values of β will be omitted from now on. The instance WS1 uses $(I_r^t)_2$; the instance WS2 uses $(I_r^t)_1$; the instances WS3 and WS4 use $(I_r^t)_1$ in R_1 , $(I_r^t)_2$ in R_3 , and $(I_r^t)_3$ in R_4 ; and the instances WS5 and WS6 use $(I_r^t)_2$ in R_1 and $(I_r^t)_3$ in R_4 .

For each reservoir $r \in R$, its capacity U_{rr} has been randomly generated as an integer in $[5000, 50,000]$, the minimum level of storage $L_{rr} = 0.01U_{rr}$, and the amount of water available at the beginning of the planning period $x_{rr}^0 = \beta(U_{rr} + L_{rr})/2$. Regarding the water users, 50% have a daily demand randomly generated in $[1, 2]$; 30% have a daily demand randomly generated in $[2, 6]$; and 20% have a daily demand randomly generated in $[6, 10]$. The daily minimum downstream requirement has been set at 10, and the maximum daily capacity of the channels has been set at 2000. Daily data are added up, as done for the inflows, depending on the periods of the planning horizon. The upper bounds U_{rw} are infinite. Table 2 displays the capacity of each reservoir, the number of demand points and of water users associated to it, and the total demand of the water users in each period of time when 12 periods are considered. Table 3 displays the total water available in each water system depending on the value of β , as well as the total demand when $T = 12$.

For each water user $k \in K_w$, $w \in W$, the economic return b_{wk}^t is randomly generated as an integer value in the interval $[10, 30]$ and it is same in all the periods. To assess the effect of varying the bounds P_{wk}^t , we have selected those bounds as a percentage of the economic return, that is, $P_{wk}^t = \alpha b_{wk}^t$, where $\alpha = 0, 0.25, 0.5, 0.75$. Note that when $\alpha = 0$ no prices are charged.

For these instances, the MLMF-BP has been solved for each combination of water system, values of α and β , and length of the planning horizon, that is, 720 problems have been solved. Moreover,

Table 2

For each reservoir: Its capacity, the demand points associated with it, the number of water users corresponding to each demand point, and the total demand of these water users in each period of time when $T = 12$

	U_{rr}	DP	$ K_w $	Total demand of each demand point in each period of time											
				1	2	3	4	5	6	7	8	9	10	11	12
R_1	35,469	1	18	1436	1523	1556	1707	1786	1856	1875	1802	1700	1584	1501	1435
R_2	38,598	2	10	979	1012	1031	1094	1121	1163	1168	1158	1090	1034	999	992
		3	5	714	721	731	795	804	826	843	794	802	743	723	712
		4	6	359	367	398	441	476	491	497	476	444	404	377	350
R_3	36,099	5	17	1629	1707	1768	1925	1985	2040	2016	1981	1892	1805	1693	1655
R_4	13,545	6	9	860	880	904	961	1017	1031	1018	1024	959	939	905	850
		7	9	866	880	937	969	1020	1031	1032	1015	977	931	868	848
R_5	42,588	8	12	1000	1049	1099	1213	1250	1282	1295	1264	1198	1110	1070	1036
		9	16	1729	1778	1865	2001	2103	2186	2162	2101	2005	1910	1782	1711
R_6	21,063	10	8	852	894	951	999	1046	1081	1098	1087	996	939	910	840
		11	15	1281	1307	1407	1472	1552	1611	1603	1553	1470	1434	1338	1284
R_7	11,760	12	17	1658	1744	1844	1993	2088	2156	2183	2110	2002	1870	1717	1670

Table 3

For each water system: The inflows involved, total water available depending on the value of β (inflow plus water available at the reservoirs at the beginning of the planning period) and total demand of its water users when $T = 12$

		β							Demand
		0.1	0.2	0.3	0.4	0.5	0.6	0.7	
WS1	$(I'_r)_2$	20,895	41,790	62,684	83,579	104,474	125,369	146,263	91,708
WS2	$(I'_r)_1$	29,558	59,117	88,675	118,234	147,792	177,350	206,909	180,945
WS3	$(I'_r)_1, (I'_r)_2, (I'_r)_3$	50,201	100,402	150,603	200,805	251,006	301,207	351,408	91,708
WS4	$(I'_r)_1, (I'_r)_2, (I'_r)_3$	53,416	106,831	160,247	213,662	267,078	320,493	373,909	157,911
WS5	$(I'_r)_2, (I'_r)_3$	32,849	65,698	98,548	131,397	164,246	197,095	229,944	128,907
WS6	$(I'_r)_2, (I'_r)_3$	34,507	69,014	103,520	138,027	172,534	207,041	241,548	180,945

the relaxed problem has been solved for each water system and value of β from 0.1 to 0.7, that is, 126 problems, and for the values of the upper level variables another 126 lower level problems have been solved to obtain the suboptimal solution.

The large number of problems solved makes it unwieldy to analyze every optimal solution obtained. To have an insight into the optimal solutions of the relaxed problem, the bilevel feasible solution associated with it (suboptimal solution) and the evolution of the optimal solution of the bilevel problem when α increases, we have selected the instance WS1 with $\beta = 0.1$ (major drought scenario) and $T = 12$. Since $f_1^1 = 0$, we pay attention to how water users are served in accordance with those solutions. We define

- K_0 as the number of water users who do not receive water in any period;
- K_1 as the number of water users who receive all their demand in all periods;
- K_{max} as the maximum, over the periods, of the water users who receive water.

Table 4

Instance WS1 with $\beta = 0.1$ and $T = 12$: Water users' satisfaction with respect to the models considered

DP	$ K_w $	Relaxed			Suboptimal			Bilevel, $\alpha = 0$			Bilevel, $\alpha = 0.25$			Bilevel, $\alpha = 0.5$		
		K_0	K_1	K_{max}	K_0	K_1	K_{max}	K_0	K_1	K_{max}	K_0	K_1	K_{max}	K_0	K_1	K_{max}
1	18	7	4	11	10	2	8	11	7	7	7	6	11	6	6	12
2	10	5	0	5	8	0	2	1	0	9	2	2	8	5	3	5
3	5	3	0	2	4	0	1	5	0	0	5	0	0	3	0	2
4	6	1	1	5	2	1	4	0	2	6	0	5	6	1	3	5
5	17	7	3	10	12	1	5	6	0	11	11	4	6	8	5	9
6	9	4	1	5	8	0	1	9	0	0	7	2	2	4	2	5
7	9	3	0	6	7	1	2	8	1	1	7	2	2	4	3	5

The first and second columns in Table 4 display the demand point and the number of its water users. The following columns show the values of K_0 , K_1 , and K_{max} for each of the above-mentioned solutions. The values corresponding to the bilevel problem with $\alpha = 0.75$ have been omitted because its optimal solution coincides with the optimal solution of the relaxed problem. Looking at these values, we can observe that the solutions are very distinct, which shows that problems (9) and (3) model very different systems, a centralized system (the relaxed problem) and a decentralized system that takes into account the reaction of the users (the bilevel problem), respectively. On the other hand, as the value of α increases the values K_0 , K_1 , and K_{max} of the bilevel solution become more and more similar to the values of the relaxed problem solution. The suboptimal solution provides the largest values of K_0 , and the lowest of K_{max} .

In the following subsections, we summarize the results of the complete experiment to give an overview of the degree to which the central authority goals are met depending on the water availability scenario. This can help decision makers to better understand the interactions among environmental requirements, demand points, water users, and prices.

4.2. Assessing the model in terms of number of variables and computing time

The aim of this section is to provide general information about the complexity of the model measured by the number of variables and constraints involved in the models and by the average computing time required for solving the problems. Table 5 displays these results. The first and second columns show the number of periods of the planning period and the water system. The third column displays the number of variables of the relaxed problem (9). The fourth and fifth columns show, respectively, the total number of variables in the bilevel model (8) and how many of them are binary variables. The sixth and seventh columns display the number of constraints of the relaxed and bilevel problems, respectively.

Finally, the two last columns display the average of the CPU time in seconds of the 7 (28) runs corresponding to the relaxed (bilevel) model. Regarding the average times, as could be expected, the relaxed model, which does not involve integer variables, is solved in very short computing times. Concerning the bilevel model, it is worth pointing out that the computing times are quite small when $T = 12$ or $T = 52$ (less than 3 and 30 seconds, respectively). Moreover, although these times increase, especially for the more complex water systems 2 and 6, when $T = 365$, reaching an aver-

Table 5

For each instance and planning period, the number of variables and constraints, and the average CPU time

		# of variables			# of constraints		Average CPU	
		Relaxed	Bilevel	(0-1)	Relaxed	Bilevel	Relaxed	Bilevel
$T = 12$	WS1	1236	4956	1860	2208	6816	0.06	0.53
	WS2	2328	9432	3552	4176	12,984	0.09	1.90
	WS3	1212	4932	1860	2184	6792	0.05	0.50
	WS4	2040	8304	3132	3672	11,436	0.09	1.03
	WS5	1668	6780	2556	3000	9336	0.08	0.82
	WS6	2352	9456	3552	4200	13,008	0.09	1.61
$T = 52$	WS1	9568	29,536	5356	21,476	8060	0.12	5.37
	WS2	18,096	56,264	10,088	40,872	15,392	0.23	20.06
	WS3	9464	29,432	5252	21,372	8060	0.13	3.77
	WS4	15,912	49,556	8840	35,984	13,572	0.19	16.86
	WS5	13,000	40,456	7228	29,380	11,076	0.17	11.72
	WS6	18,200	56,368	10,192	40,976	15,392	0.22	27.64
$T = 365$	WS1	37,595	150,745	56,575	67,160	207,320	1.49	246.23
	WS2	70,810	286,890	108,040	127,020	394,930	3.47	1734.69
	WS3	36,865	150,015	56,575	66,430	206,590	2.38	268.60
	WS4	62,050	252,580	95,265	111,690	347,845	2.75	778.01
	WS5	50,735	206,225	77,745	91,250	283,970	1.76	895.76
	WS6	71,540	287,620	108,040	127,750	395,660	2.93	1346.72

age computing time of almost 30 minutes, these times can be considered competitive taking into account both the number of variables and constraints involved. The detailed results regarding the computing time are presented in the Appendix. It is worth mentioning at this point that, even if the bilevel problem is much more complicated to solve, it is the appropriate model when modeling a decentralized system that takes into account the reaction of the users, as noted when describing the MLMF-BP. In the previous and following sections, we show that the suboptimal solution is not necessarily optimal for the MLMF-BP. In fact, it can be very far from the bilevel optimal solution, and *a priori* it is not possible to know what extent.

4.3. Degree to which the goals of the central authority are met

For all the instances and available water, it is possible to satisfy environmental and sustainability issues, that is, the minimum needs are always met and thus f_1^1 is always equal to zero. For small values of β it is not possible to satisfy the demand of all the water users, even if the reaction of the demand points is not taken into account. The level of satisfaction of demand users, and so the second goal of the central authority, depends largely on the water system analyzed. To satisfy all the water users in all periods requires $\beta \geq 0.3$ for instance WS3 and $\beta \geq 0.4$ for instance WS4, which have a similar structure with a river and two tributaries. Instances WS1 and WS5 require $\beta \geq 0.5$. Finally, the most demanding systems from the point of view of the computing time required are also the instances that need a larger value of β to provide total satisfaction. These are WS6 that requires $\beta \geq 0.6$ and WS2 that requires $\beta \geq 0.7$. In these cases, the demand of all the

Table 6

The satisfaction index I for $T = 12$

	β	Relaxed		Bilevel			
		Optimal	Suboptimal	$\alpha = 0$	$\alpha = 0.25$	$\alpha = 0.5$	$\alpha = 0.75$
WS1	0.1	0.42	0.19	0.27	0.35	0.41	=
	0.2	0.69	0.43	0.53	0.60	0.66	=
	0.3	0.85	0.69	0.75	0.79	0.83	=
	0.4	0.94	0.87	0.92	0.93	=	=
WS2	0.1	0.33	0.16	0.22	0.30	=	=
	0.2	0.59	0.31	0.40	0.49	0.56	=
	0.3	0.73	0.49	0.57	0.64	0.70	=
	0.4	0.84	0.68	0.73	0.78	0.82	=
	0.5	0.91	0.81	0.87	0.89	=	=
	0.6	0.97	0.94	0.96	=	=	=
WS3	0.1	0.72	0.49	0.57	0.63	0.69	=
	0.2	0.95	0.88	0.92	0.93	=	=
WS4	0.1	0.59	0.33	0.41	0.49	0.56	=
	0.2	0.84	0.68	0.74	0.78	0.82	=
	0.3	0.96	0.92	0.95	=	=	=
WS5	0.1	0.44	0.19	0.27	0.35	0.43	=
	0.2	0.72	0.46	0.55	0.62	0.68	=
	0.3	0.87	0.72	0.78	0.82	0.86	=
	0.4	0.96	0.91	0.95	=	=	=
WS6	0.1	0.35	0.16	0.23	0.31	=	
	0.2	0.62	0.35	0.44	0.52	0.59	=
	0.3	0.78	0.56	0.63	0.69	0.75	=
	0.4	0.88	0.75	0.80	0.84	0.87	=
	0.5	0.95	0.89	0.93	=	=	=

water users is satisfied without the need of charging a price to the demand points. Needless to say, when there is “enough” water there are no conflicts. It is not important whether it is the central authority or the demand points that take control of the allocation of water to water users because all the objective functions achieve their best theoretical values ($f_1^1 = 0$, $f_1^2 = |T| \times |K|$, $f_1^3 = 0$, $f_2 = \sum_{t \in T} \sum_{w \in W} \sum_{k \in K_w} b_{wk}^t D_{wk}^t$). In the remaining cases in which there is not enough water to satisfy all the demands the relevance of the bilevel model and the importance of prices is made clear. As can be expected, the more water available (β is larger), the more water users have their demand satisfied.

In what follows, we consider the values of β for which it is not possible to satisfy the total demand without charging prices, and pay attention to the satisfaction of water users measured by the so-called satisfaction index defined as the quotient between the total satisfaction given by the value of f_1^2 evaluated for the corresponding solution and the satisfaction of all the water users:

$$I = \frac{f_1^2}{|T| \times |K|}.$$

Table 6 summarizes the results for $T = 12$. As mentioned above, this table does not include those values of β for which this index is equal to 1, that is, every water user receives as much water as he/she demands in every period. The first and the second columns of the table display the water system and the value of β . The third and fourth columns show the satisfaction index I value of the optimal solution of the relaxed problem and the corresponding suboptimal solution, respectively. Finally, the remaining columns display the value of I associated with the optimal solution of the bilevel problem when α is equal to 0, 0.25, 0.5, and 0.75. A symbol “=” is written when the index value coincides with the index value provided by the optimal solution of the relaxed problem. This value is an upper bound of the index value of every bilevel feasible solution. The results corresponding to $T = 52$ and $T = 365$ are shown in the Appendix. Since the results in all the tables are very similar (they are better for $T = 365$ than for $T = 52$, which is better than for $T = 12$, but the differences are less than 0.02), we summarize the information for $T = 12$.

Looking at Table 6, as could be expected, the satisfaction index associated with every instance always increases as the available water is increased (β grows). The optimal solution of the relaxed problem provides the best values but, as mentioned above, in a hierarchical framework, the suboptimal solution is obtained and thus the satisfaction index actually computed is the one shown under suboptimal, which always provides the worst values. However, comparing the index value of the relaxed optimal solution with the indexes of the optimal solution of the bilevel problems, we can realize which value of α provides the best satisfaction and which value provides an index close enough to that. Note also that to obtain the optimal solution of the relaxed problem by means of a bilevel problem, in general it is needed to enforce $\alpha = 0.75$, the largest value analyzed in this study. For large enough values of β , this also happens when $\alpha = 0.5$ and $\alpha = 0.25$ for some water systems.

5. Conclusions

In this paper we have proposed a decision tool based on bilevel optimization, which can be used by decision makers to be aware of the impact of different water-allocation policies that can involve charging fees when it is mandatory to take into account the decisions made at the next level of the decision-making process. For this purpose, we have proposed a multiobjective multifollower bilevel optimization model to manage water allocation in a hierarchical decentralized water system. The purpose is to guarantee environmental requirements as well as to make an efficient use of available water in terms of satisfying demand. Due to the hierarchical decision process involved, a bilevel model is proposed whose upper level decision maker, the central authority, aims first to satisfy minimum requirements. Second, he/she aims to maximize the overall satisfaction of the water user demands. In the process of deciding how to allocate the available water, the upper level is constrained by the behavior of the demand point managers who are at the lower level of the decision process. At this level, there are as many decision makers as demand points, each of them aiming to distribute the water assigned to them in accordance with the economic return. In addition, the central authority also has as a regulatory tool the possibility of setting fees to be paid by the demand point depending on how it decides to allocate the water. The model allows us to see the evolution of the users satisfaction as the value of prices charged and the water available vary. The model can be easily adapted to manage very different water systems as can be inferred from the very distinct water systems and drought scenarios dealt with in the computational experiments.

The model resulting from this approach, a multiobjective lexicographic linear bilevel optimization problem, has been reformulated as a multiobjective lexicographic linear bilevel problem with one follower. Next, an exact procedure to solve the model has been proposed based on applying the KKT conditions of the lower level problem. Valid constants for the big- M s involved in the linearization of the KKT conditions are derived.

The results of the extensive computational experiments carried out on a large set of benchmark instances, which considers different water systems, drought scenarios, and several planning periods, confirm the efficiency of the procedure since it is able to provide the optimal solution in reasonable computing times (very short in most cases). This avoids the need to develop heuristics or meta-heuristics algorithms since commercial software can be used to solve the problems involved. Moreover, it allows us to foresee that this procedure will be key when dealing with stochastic inflows or demands in future research.

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Appendix

Tables A1–A6, respectively, display the computing times required for solving the MLMF-BP and the relaxed problem for the instances WS1–WS6. In each table, the first and second columns display, respectively, the value of the coefficients β and α . The remaining six columns display t_{bil} , the CPU time in seconds required to solve the MLMF-BP, and t_{rel} , the CPU time in seconds required by CPLEX to solve the relaxed problem depending on the periods of the planning horizon, respectively. The CPLEX stopping criterion was set at 7200 seconds. Only 5 of the 720 bilevel problems (all of them having $T = 365$) were stopped before providing an optimal solution.

Table A1
Computing times of the instance WS1

β	α	$T = 12$		$T = 52$		$T = 365$	
		t_{bil}	t_{rel}	t_{bil}	t_{rel}	t_{bil}	t_{rel}
0.1	0	1.07	0.04	13.89	0.11	180.94	3.55
	0.25	1.67		20.22		126.77	
	0.5	0.49		8.54		153.42	
	0.75	0.33		1.82		70.69	
0.2	0	0.96	0.04	5.35	0.10	7.75	1.36
	0.25	1.44		34.63		7.45	
	0.5	0.90		8.66		7.59	
	0.75	0.66		3.24		561.90	
0.3	0	0.89	0.05	9.57	0.11	2221.31	1.35
	0.25	1.47		12.84		1285.28	
	0.5	0.85		7.63		96.73	
	0.75	0.30		2.09		7.34	
0.4	0	0.60	0.05	4.34	0.11	7.23	1.37
	0.25	0.63		5.04		7.23	
	0.5	0.68		4.48		115.26	
	0.75	0.29		1.84		1057.37	
0.5	0	0.15	0.06	0.58	0.10	607.91	0.94
	0.25	0.13		0.56		62.66	
	0.5	0.12		0.50		5.98	
	0.75	0.11		0.46		5.91	
0.6	0	0.15	0.05	0.59	0.10	5.92	0.93
	0.25	0.13		0.55		31.06	
	0.5	0.12		0.49		57.63	
	0.75	0.11		0.46		160.39	
0.7	0	0.15	0.06	0.59	0.10	26.25	0.93
	0.25	0.13		0.55		5.55	
	0.5	0.12		0.49		5.44	
	0.75	0.11		0.46		5.42	

Table A2
Computing times of the instance WS2

β	α	$T = 12$		$T = 52$		$T = 365$	
		t_{bil}	t_{rel}	t_{bil}	t_{rel}	t_{bil}	t_{rel}
0.1	0	3.32	0.08	31.73	0.20	559.17	5.14
	0.25	3.50		25.26		296.83	
	0.5	1.91		11.82		582.40	
	0.75	0.57		6.36		583.45	
0.2	0	3.27	0.08	21.65	0.22	409.48	3.31
	0.25	4.50		28.48		134.16	
	0.5	2.28		52.09		16.43	
	0.75	0.54		6.91		2081.76	
0.3	0	2.39	0.08	25.87	0.22	7200.66	3.40
	0.25	4.40		36.27		7200.65	
	0.5	1.27		49.46		6849.79	
	0.75	0.59		9.97		3747.03	
0.4	0	2.16	0.08	27.83	0.22	269.39	3.38
	0.25	3.75		108.57		15.83	
	0.5	3.47		15.50		427.58	
	0.75	0.54		3.69		6428.74	
0.5	0	1.99	0.08	19.18	0.21	3370.57	3.40
	0.25	3.26		25.89		3927.89	
	0.5	1.43		18.14		373.39	
	0.75	0.77		3.64		175.73	
0.6	0	1.33	0.08	9.43	0.22	13.52	3.55
	0.25	2.39		10.31		945.41	
	0.5	1.57		6.08		183.59	
	0.75	0.70		2.86		902.10	
0.7	0	0.28	0.07	1.29	0.18	1617.57	2.13
	0.25	0.47		1.34		134.22	
	0.5	0.49		1.08		111.55	
	0.75	0.20		0.99		12.41	

Table A3
Computing times of the instance WS3

β	α	$T = 12$		$T = 52$		$T = 365$	
		t_{bil}	t_{rel}	t_{bil}	t_{rel}	t_{bil}	t_{rel}
0.1	0	2.51	0.04	11.68	0.11	222.34	3.41
	0.25	3.40		56.64		36.52	
	0.5	1.20		8.01		9.54	
	0.75	0.95		1.87		10.25	
0.2	0	0.42	0.04	2.88	0.10	10.25	1.95
	0.25	0.71		4.27		10.25	
	0.5	0.49		4.64		10.37	
	0.75	0.47		1.46		5417.98	

Continued

Table A3
(Continued)

β	α	$T = 12$		$T = 52$		$T = 365$	
		t_{bil}	t_{rel}	t_{bil}	t_{rel}	t_{bil}	t_{rel}
0.3	0	0.17	0.04	0.70	0.10	158.81	3.85
	0.25	0.19		0.69		11.06	
	0.5	0.17		0.61		11.58	
	0.75	0.16		0.58		11.55	
0.4	0	0.20	0.04	0.73	0.11	11.48	3.12
	0.25	0.19		0.82		11.48	
	0.5	0.17		0.73		1213.20	
	0.75	0.17		0.67		63.62	
0.5	0	0.19	0.04	0.72	0.10	10.02	1.80
	0.25	0.19		0.80		10.34	
	0.5	0.17		0.69		9.87	
	0.75	0.16		0.61		10.22	
0.6	0	0.59	0.04	0.75	0.11	9.80	1.32
	0.25	0.20		0.83		175.88	
	0.5	0.17		0.72		27.68	
	0.75	0.17		0.61		9.42	
0.7	0	0.18	0.04	0.74	0.11	9.82	1.25
	0.25	0.19		0.93		9.23	
	0.5	0.18		0.72		9.12	
	0.75	0.16		0.61		9.13	

Table A4
Computing times of the instance WS4

β	α	$T = 12$		$T = 52$		$T = 365$	
		t_{bil}	t_{rel}	t_{bil}	t_{rel}	t_{bil}	t_{rel}
0.1	0	1.78	0.07	20.61	0.20	193.90	6.74
	0.25	8.82		255.16		323.54	
	0.5	2.30		46.31		149.58	
	0.75	0.81		3.31		13.61	
0.2	0	1.54	0.08	12.88	0.18	13.70	2.88
	0.25	2.84		64.38		13.76	
	0.5	1.43		17.50		13.77	
	0.75	0.48		5.42		7201.24	
0.3	0	1.12	0.08	8.05	0.20	5235.90	2.74
	0.25	1.90		11.37		534.20	
	0.5	1.41		7.92		14.94	
	0.75	0.45		3.30		14.99	
0.4	0	0.23	0.07	1.00	0.16	15.10	1.76
	0.25	0.26		1.04		15.07	
	0.5	0.21		1.00		5064.09	
	0.75	0.19		0.93		1615.83	

Continued

Table A4
(Continued)

β	α	$T = 12$		$T = 52$		$T = 365$	
		t_{bil}	t_{rel}	t_{bil}	t_{rel}	t_{bil}	t_{rel}
0.5	0	0.23	0.07	1.01	0.16	155.96	1.75
	0.25	0.26		1.06		12.60	
	0.5	0.63		0.88		12.29	
	0.75	0.18		0.83		12.27	
0.6	0	0.23	0.07	1.01	0.15	12.15	1.72
	0.25	0.26		1.06		652.18	
	0.5	0.21		0.89		374.32	
	0.75	0.18		1.01		73.87	
0.7	0	0.23	0.07	1.01	0.16	11.68	1.68
	0.25	0.26		1.48		11.27	
	0.5	0.22		0.89		11.22	
	0.75	0.18		0.82		11.14	

Table A5
Computing times of the instance WS5

β	α	$T = 12$		$T = 52$		$T = 365$	
		t_{bil}	t_{rel}	t_{bil}	t_{rel}	t_{bil}	t_{rel}
0.1	0	1.52	0.05	43.88	0.15	276.72	2.24
	0.25	3.46		105.33		260.61	
	0.5	1.40		9.63		271.90	
	0.75	0.46		2.35		115.36	
0.2	0	1.48	0.07	16.85	0.14	10.42	2.03
	0.25	2.71		50.90		10.47	
	0.5	1.31		14.84		10.57	
	0.75	0.41		2.77		4544.62	
0.3	0	1.40	0.07	10.19	0.14	3892.10	2.05
	0.25	2.09		24.26		5908.68	
	0.5	1.06		10.43		227.78	
	0.75	0.35		3.75		11.88	
0.4	0	0.86	0.07	5.86	0.13	11.88	2.08
	0.25	1.35		9.05		11.89	
	0.5	0.70		5.30		2817.52	
	0.75	0.38		3.44		4995.98	
0.5	0	0.19	0.07	0.74	0.12	754.21	1.36
	0.25	0.20		0.78		121.10	
	0.5	0.16		0.81		9.80	
	0.75	0.15		0.82		9.64	
0.6	0	0.19	0.07	0.72	0.13	9.63	1.30
	0.25	0.20		0.78		82.20	
	0.5	0.17		0.80		438.61	
	0.75	0.15		0.73		205.93	

Continued

Table A5
(Continued)

β	α	$T = 12$		$T = 52$		$T = 365$	
		t_{bil}	t_{rel}	t_{bil}	t_{rel}	t_{bil}	t_{rel}
0.7	0	0.19	0.07	0.72	0.12	45.22	1.27
	0.25	0.20		0.78		8.86	
	0.5	0.17		0.78		8.84	
	0.75	0.15		0.80		8.80	

Table A6
Computing times of the instance WS6

β	α	$T = 12$		$T = 52$		$T = 365$	
		t_{bil}	t_{rel}	t_{bil}	t_{rel}	t_{bil}	t_{rel}
0.1	0	3.34	0.08	70.59	0.24	485.46	3.47
	0.25	3.75		99.56		346.70	
	0.5	1.50		12.81		629.89	
	0.75	0.63		3.15		534.05	
0.2	0	1.30	0.08	16.04	0.23	242.78	3.20
	0.25	3.94		24.02		16.18	
	0.5	1.51		55.97		16.03	
	0.75	1.11		4.90		2264.29	
0.3	0	2.61	0.08	19.23	0.22	7200.61	3.39
	0.25	4.99		244.81		7200.62	
	0.5	3.19		65.83		4041.26	
	0.75	0.64		3.72		357.03	
0.4	0	2.26	0.08	15.29	0.22	15.77	3.33
	0.25	2.86		69.67		15.62	
	0.5	1.82		10.72		1938.42	
	0.75	0.79		3.88		5105.76	
0.5	0	1.33	0.08	19.99	0.22	3743.76	3.23
	0.25	2.11		13.90		1225.49	
	0.5	2.26		8.38		259.25	
	0.75	0.49		2.53		13.30	
0.6	0	0.25	0.08	1.27	0.19	13.27	1.97
	0.25	0.24		1.34		357.14	
	0.5	0.22		0.95		94.36	
	0.75	0.19		0.89		395.56	
0.7	0	0.66	0.08	1.26	0.19	743.10	1.93
	0.25	0.24		1.33		428.40	
	0.5	0.60		0.95		12.14	
	0.75	0.19		0.93		12.05	

Tables A7 and A8 display the satisfaction index for $T = 12$ and $T = 365$, respectively. They have the same structure as Table 6.

Table A7

The satisfaction index I for $T = 52$

	β	Relaxed		Bilevel			
		Optimal	Suboptimal	$\alpha = 0$	$\alpha = 0.25$	$\alpha = 0.5$	$\alpha = 0.75$
WS1	0.1	0.42	0.19	0.28	0.35	=	=
	0.2	0.70	0.43	0.53	0.60	0.67	=
	0.3	0.85	0.69	0.75	0.80	0.83	=
	0.4	0.94	0.87	0.92	0.93	=	=
WS2	0.1	0.34	0.16	0.22	0.30	0.33	=
	0.2	0.60	0.32	0.41	0.49	0.56	=
	0.3	0.74	0.49	0.58	0.64	0.70	=
	0.4	0.84	0.68	0.73	0.78	0.82	=
	0.5	0.91	0.82	0.87	0.89	=	=
	0.6	0.97	0.94	0.96	=	=	=
WS3	0.1	0.72	0.50	0.58	0.64	0.70	=
	0.2	0.95	0.88	0.92	0.93	=	=
WS4	0.1	0.60	0.33	0.42	0.50	0.57	=
	0.2	0.84	0.68	0.74	0.79	0.83	=
	0.3	0.96	0.92	0.95	=	=	=
WS5	0.1	0.45	0.20	0.27	0.36	0.43	=
	0.2	0.72	0.47	0.55	0.62	0.69	=
	0.3	0.87	0.73	0.79	0.82	0.86	=
	0.4	0.96	0.91	0.95	=	=	=
WS6	0.1	0.36	0.16	0.23	0.31	0.35	=
	0.2	0.63	0.35	0.44	0.52	0.59	=
	0.3	0.78	0.57	0.64	0.70	0.75	=
	0.4	0.88	0.76	0.81	0.84	0.87	=
	0.5	0.95	0.90	0.93	=	=	=

Table A8

The satisfaction index I for $T = 365$

	β	Relaxed		Bilevel			
		Optimal	Suboptimal	$\alpha = 0$	$\alpha = 0.25$	$\alpha = 0.5$	$\alpha = 0.75$
WS1	0.1	0.43	0.20	0.28	0.36	=	=
	0.2	0.71	0.43	0.54	0.61	0.68	=
	0.3	0.85	0.69	0.76	0.80	0.84	=
	0.4	0.94	0.87	0.92	=	=	=
WS2	0.1	0.35	0.16	0.22	0.30	0.34	=
	0.2	0.61	0.32	0.41	0.50	0.58	=

Continued

Table A8
(Continued)

	β	Relaxed		Bilevel			
		Optimal	Suboptimal	$\alpha = 0$	$\alpha = 0.25$	$\alpha = 0.5$	$\alpha = 0.75$
	0.3	0.74	0.49	0.58	0.65	0.72	=
	0.4	0.84	0.68	0.74	0.79	0.82	=
	0.5	0.92	0.82	0.87	0.90	0.91	=
	0.6	0.97	0.94	=	=	=	=
WS3	0.1	0.73	0.50	0.59	0.65	0.71	=
	0.2	0.95	0.89	0.92	0.94	=	=
WS4	0.1	0.61	0.33	0.42	0.51	0.58	=
	0.2	0.85	0.68	0.75	0.79	0.83	=
	0.3	0.96	0.92	0.95	=	=	=
WS5	0.1	0.45	0.20	0.27	0.36	0.43	=
	0.2	0.73	0.47	0.56	0.63	0.70	=
	0.3	0.88	0.73	0.79	0.83	0.86	=
	0.4	0.96	0.91	0.95	=	=	=
WS6	0.1	0.37	0.17	0.23	0.31	0.36	=
	0.2	0.65	0.36	0.45	0.54	0.61	=
	0.3	0.78	0.57	0.65	0.71	0.76	=
	0.4	0.88	0.76	0.81	0.85	0.87	=
	0.5	0.95	0.90	0.94	=	=	=